



Most-Probable: A New Argumentation Semantics through Optimization



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Introduction

> Argumentation has origin from philosophy, but has gathered interests in many fields (e.g. law).

- ➢ We refer to the Dung's formalization [2,3].
- Extensions to represent uncertainty are available.
- Representation techniques are often borrowed from the KRR field.



[2] P. M. Dung, On the acceptability of arguments and its fundamental role in nonmonotonic reasoning, logic programming and n-person games, Artificial intelligence 77 (1995) 321-357.

[3] A. Bondarenko, P. M. Dung, R. A. Kowalski, F. Toni, An abstract, argumentation-theoretic approach to default reasoning, Artificial intelligence 93 (1997) 63-101.

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Answer Set Programming (ASP)

An Answer Set Program is a program made up of fact(s) like

 $p(t_1, ..., t_p).$

and rule(s) like

$$H := A_1, \dots, A_n, \text{ not } B_1, \dots, \text{ not } B_{m.}$$

Special rule(s) have no head (constraints).

$$:-A_1, ..., A_n, \neg B_1, ..., \neg B_m.$$

ASP supports aggregates (sum, count, ...) and optimizations (optimize).

Objectives

We propose a strategy for updating beliefs in an argumentation network based on attack relationships.

We propose a new semantics (*most-probable*) that deals with degrees of belief.

This semantics focuses on a *target* argument, trying to collect all the arguments that justify (or defend) the target maximizing degrees of belief.

We integrated the belief updating and the semantics computation with an existing Prolog-based system for argumentation.

Background

➤ **Definition** An Argumentation Framework (AF) is a pair $\langle Args, Att \rangle$ where Args is a (finite) set of arguments and $Att \subseteq Args \times Args$ the attack relationship. The concept of attack can also be extended to a set of arguments. Given $S \subseteq Args, A \in Args$ attacks S iff exists $B \in S$ such that A attacks B.



Background

▶ **Definition** Let an $AF \ G = \langle Args, Att \rangle$ be an AF, a set of arguments $S \subseteq Args$ is conflict-free iff $\langle A, B \rangle \in S$ such that $\langle A, B \rangle \in Att$. The set of all conflict-free extensions is indicated as cf(G).





Background

➤ Definition Let an AF G = $\langle Args, Att \rangle$ be an AF, a set of arguments S ⊆ Args is admissible iff S is conflict-free and ∀γ ∈ Args γ attacks S ⇒ S attacks γ. In other words, S defends every A ∈ S.





Simplified Probabilistic Argumentation Framework

- Probabilities come under the assumption of non-additivity, meaning that it is not impossible that $\Delta \nvDash \alpha$ and $\Delta \nvDash \neg \alpha$.
- Attacks always trigger with the same intensity.
- Probabilities of arguments cannot be compared with Kolmogorov Axioms.
- ➤ Definition Let an AF G = $\langle Args, Att \rangle$, a Simplified Probabilistic Argumentation Framework (SPAF) $\langle Args, Att, PArgs \rangle$ is a triple in which Args and Att are defined as above, PArgs : Args →]0, 1], a function indicating the likelihood of arguments.

$$A \longrightarrow B \longrightarrow C \longrightarrow D$$

$$P(A) = 0.8 \quad P(B) = 0.4 \quad P(C) = 0.7 \quad P(D) = 0.3$$

- Probabilities of arguments are "blind" with respect to other arguments and attacks.
- We update the probabilities of arguments, based on attacks and the initial probability of arguments.
- The proposed formula is

$$P'(A) = P(A) \cdot \prod_{\substack{\gamma \in Args \\ \langle \gamma, A \rangle \in Att}} 1 - \alpha \cdot P(\gamma)$$

where P'(A) represents the probability of A after being attacked and $\alpha \in [0, 1]$.

The nice properties are:

- multiple attacks influence much the updated belief.
- $\blacktriangleright \quad \text{the product } 1 \alpha \cdot P(\gamma) \text{ is always between 0 and 1.}$
- the initial belief of an argument is an upper bound for the update.
- > many weak (low $P(\gamma)$) attackers are less influential than few strong (high $P(\gamma)$) attackers.

$$P'(A) = P(A) \cdot \prod_{\substack{\gamma \in Args \\ \langle \gamma, A \rangle \in Att}} 1 - \alpha \cdot P(\gamma)$$

The main limitations are:

- ➢ if an attacker has probability 1, the attacked is not vanished.
- the function is nonlinear, meaning that small changes have large effects.
- ➢ it is assumed attacks are independent.

$$P'(A) = P(A) \cdot \prod_{\substack{\gamma \in Args \\ \langle \gamma, A \rangle \in Att}} 1 - \alpha \cdot P(\gamma)$$



We may also consider an iterative updating probability step.

$$P^{n}(A) = \begin{cases} P(A) & \text{if } n = 0, \\ P^{n-1}(A) \cdot U^{n}(A) & otherwise \end{cases}$$

where

$$U^{i}(A) = P^{i-1}(A) \cdot \prod_{\substack{\gamma \in Args \\ \langle \gamma, A \rangle \in Att}} 1 - \alpha \cdot P^{i-1}(\gamma).$$

The proposed formula resembles one of those proposed by Gabbay et al. [11]. However, the starting probabilities of the arguments were not taken into account.

^[11] D. M. Gabbay, Equational approach to argumentation networks, Argument & Computation 3 (2012) 87–142.

Most-Probable Semantics

- Definition Let (Args, Att, P Args) G a Simplified Probabilistic Argumentation Framework and t ∈ Args, a set of arguments S ⊆ Args is a most-probable extension for t, indicated as S ∈ mostprob_t(G) iff t ∈ S, S is admissible and $\exists Y \subset S$ such that:
 - $\succ t \in Y$
 - ➤ Y is admissible
 - $\succ P(Y \setminus \{t\}) > P(S \setminus \{t\}).$
 - and $\nexists Z \subseteq Args$ such that:
 - $\succ S \subset Z$
 - \succ *Z* is admissible
 - $\succ P(Z \setminus \{t\}) = P(S \setminus \{t\}).$

$$P(\{a_1, ..., a_n\}) = \min\{P(a_1), ..., P(a_n)\}$$

Most-Probable Semantics

Lemma There is not a unique solution for the most-probable extension.

Proof. By construction, suppose *G* a SPAF with $Args = \{a_1, a_2, ..., a_n, s_0, s_1, t\}$ where $\{a_1, a_2, ..., a_n, t\}$ are not attacked by anyone, and s_0, s_1 attack each other. Then, both $\{a_1, a_2, ..., a_n, s_0, t\}$ and $\{a_1, a_2, ..., a_n, s_1, t\}$ are most-prob_t(*G*).



δ -most-Probable Semantics

- ➤ Definition Let (Args, Att, P Args) G a Simplified Probabilistic Argumentation Framework, $t \in Args$ and $\delta \in \mathbb{R}_>$, a set of arguments $S \subseteq Args$ is a δ-most-probable extension with respect to t iff S is most-probable with respect to t and $P(S \setminus \{t\}) >= \delta$.
 - $\delta = 0$ would reduce threshold most-probable to most-probable.

Implementation

We developed this probabilistic argumentation framework in a platform for argument reasoning called ARGuing Using Enhanced Reasoning (Arguer).



Implementation

- Arguer provides the following argumentation frameworks:
 - Abstract Argumentation Framework (AF)
 - Value-Based AF
 - Bipolar AF
 - Weighted AF
 - Bipolar Weighted AF
 - Simplified Probabilistic AF

Implementation (Probability Update)

```
update(Arguments, Attacks, Likelihoods) :-
    adjust likelihoods(Arguments, Attacks, Likelihoods), normalize probabilities.
adjust likelihoods([], , ).
adjust likelihoods([Arg | Rest], Attacks, Likelihoods) :- adjust likelihood(Arg, Attacks, Likelihoods),
    adjust likelihoods(Rest, Attacks, Likelihoods).
adjust likelihood(Arg, Attacks, Likelihoods) :-
    findall(X, lists:member([X,Arg], Attacks), Attackers),
    findall(X, (lists:member(A, Attackers), lists:member([A,X], Likelihoods)), ProbabilitiesAttackers),
    lists:member([Arg, OriginalProbability], Likelihoods), Alfa is 0.2,
    compute product(Alfa, ProbabilitiesAttackers, Product), AdjustedProbability is Product * OriginalProbability,
    assertz(not normalized probability(Arg, AdjustedProbability)).
compute product(Alfa, [], 1) :-
.
compute product(Alfa, [P | Rest], Product) :-
    TermProduct is 1 - Alfa * P, compute product(Alfa, Rest, OldProduct), Product is TermProduct * OldProduct.
find max([X], X).
find max([Number | Rest], Max) :-
    find max(Rest, MaxRest), Max is max(Number, MaxRest).
normalize probabilities :-
    findall(X, not normalized probability(,X), Probabilities), find max(Probabilities, Max),
    Factor is 100 / Max, findall((Arg, Probability),
    not normalized probability(Arg, Probability), ArgProbabilities), assert probabilities(ArgProbabilities, Factor).
assert probabilities([], Factor) :-
1.
assert probabilities([ArgProbability | Rest], Factor) :-
    ArgProbability = (Argument, Probability), NormalizedProbability is round(Factor * Probability),
    assertz(probability(Argument, NormalizedProbability)), assert probabilities(Rest, Factor).
```

Implementation (Most-Probable)

```
argument(a;b;c;d;e;f).
target(d).
attack(a, b). attack(a, b).
attack(b, c). attack(b, d).
attack(c, a). attack(e, b).
attack(f, b).
likelihood(a, 80).
likelihood(b, 40).
likelihood(c, 80).
likelihood(d, 30).
likelihood(e, 70).
likelihood(f, 40).
% the target is always in the most-prob
most prob(N) :- target(N).
%all the others are candidate
{most prob(N) : argument(N)} :- not target(N).
%arguments must be defended (or attacks free) to be good candidates
:- most prob(A), attack(B, A), not most prob(C) : attack(C, B).
%take minimum likelihood
min likelihood(Min) :- Min = #min { L : most prob(A), likelihood(A, L), not target(A) }.
%take cardinality
size most prob(C) :- C = #count { A : most prob(A) }.
%take set with the maximum minimum likelihood
#maximize { L : min likelihood(L) }.
%add as many arguments as possible without affecting the minimum likelihood
#maximize { C : size most prob(C) }.
#show most prob/1.
```

Conclusions

- We proposed:
 - a new representation for probabilistic argumentation frameworks.
 - a new semantics, with implementation and interpretation of it.
- Much research can be pursued in:
 - mitigating the current limitations of probability updates.
 - > proposing new interpretations for initial probabilities.
 - suitable applications.